Turning Failure into Proof: The ProB Disprover

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Abstract

Initially, the ProB disprover used constraint solving to try and find counterexamples to proof obligations generated from Event-B models [1]. Recently, we made the ProB kernel capable of determining whether a search was exhaustive. Hence, one can now also use it as a prover. It uses a technique which guarantees soundness in the following way: if a solution is found it is guaranteed to be correct; if not, the solver can either return the result “definitely false” if the enumeration was exhaustive, or “unknown” if the enumeration was aborted.

Technique

ProB’s constraint solver is based on CLP(FD)-style constraint propagation and resorts to enumeration once no further propagation is possible. Hence, it relies on variables being constrained to a finite domain. While this is fine for animation and visualization, it limits the applicability of ProB as a prover. To overcome this limitation, we decided to track the enumeration of certain variables. We firstly set up an environment for each scope (i.e. quantifier, set comprehension, and so on), stating whether a variable has to be enumerated exhaustively for a proof to be successful. Then, we use this information to direct the enumeration towards certain variables, knowing whether a solution means satisfiability or unsatisfiability or if further search is necessary. With this technique, we are able to prove if certain hypotheses $H$ imply a Goal by searching for a solution to $H \implies \neg \text{Goal}$.

Results

We applied the ProB disprover to Event-B proof obligations [2] and compared it to the classical B provers as well as to the SMT solver based prover. Additionally, we wrote a translation from SMT-LIB to B and used problems selected from the SMT-LIB collection of benchmarks. Here, we compared to CVC4 and Z3. The diagrams below show the number of discharged / solved problems.

Conclusion

As explained before, ProB’s technique differs from classic DPLL(T)-based provers, making it an orthogonal addition to them rather than a replacement. Our empirical evaluation shows, that our approach is useful for certain kinds of problems which are otherwise hard to solve. We therefore think that it should be evaluated and refined further.

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